Further investigations on the QAM method for finding new APN functions

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- \mathbb{F}_{2^n} finite field with 2^n elements, $n \in \mathbb{N}$.
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- ► $\delta_F = \max_{a,b \in \mathbb{F}_{2^n}, a \neq 0} |\{x \in \mathbb{F}_{2^n} : \Delta_F(a,x) = b\}|$ its differential unifomity.



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- *F* is almost perfect nonlinear(APN) if $\delta_F = 2$.



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Preliminaries 00●0000	Matrix structure 00	Orbit restrictions	Submatrix method 00	Computational searches

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- *F* is a **linear** function if $F(x) = \sum_{0 \le i < n} a_i x^{2^i}$, $a_i \in \mathbb{F}_{2^n}$.

F is **affine** if it is a sum of a linear and a constant.



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Preliminaries 00●0000	Matrix structure 00	Orbit restrictions	Submatrix method 00	Computational searches

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- F is quadratic if $deg(F) \leq 2$.
- We will consider homogeneous quadratic (n, n)-function F

$$F(x) = \sum_{0 \le i < j \le n-1} a_{i,j} x^{2^i + 2^j}, \ a_{i,j} \in \mathbb{F}_{2^n}.$$



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The functions F and F' from \mathbb{F}_{2^n} to itself are called

■ affine equivalent (or linear equivalent) if F' = A₁ ∘ F ∘ A₂ for affine (linear) permutations A₁, A₂ from F_{2ⁿ} to itself.



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 For quadratic APN (n, n) functions, F and F' are CCZ-equivalent if and only if they are EA-equivalent [2].



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- [3] The rank of the vector v ∈ Fⁿ_{2ⁿ} is the dimension of the subspace spanned by its elements.



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- [3] The rank of the vector v ∈ ℝⁿ_{2ⁿ} is the dimension of the subspace spanned by its elements.
- ▶ The **derivative matrix** $M_F \in \mathbb{F}_{2^n}^{n \times n}$ of function *F* is

$$M_{F} = \begin{bmatrix} \Delta_{b}F(b) & \Delta_{b}F(b^{2}) & \dots & \Delta_{b}F(b^{n}) \\ \Delta_{b^{2}}F(b) & \Delta_{b^{2}}F(b^{2}) & \dots & \Delta_{b^{2}}F(b^{n}) \\ \vdots & \vdots & \ddots & \vdots \\ \Delta_{b^{n}}F(b) & \Delta_{b^{n}}F(b^{2}) & \dots & \Delta_{b^{n}}F(b^{n}) \end{bmatrix}$$



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A matrix M_F ∈ ℝ^{n×n}_{2ⁿ} is called a Quadratic APN Matrix (QAM) [3] if:



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 - 1. M_F is symmetric and the elements in its main diagonal are all zeros;



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- A matrix M_F ∈ ℝ^{n×n}_{2ⁿ} is called a Quadratic APN Matrix (QAM) [3] if:
 - 1. M_F is symmetric and the elements in its main diagonal are all zeros;
 - 2. Every nonzero linear combination of the *n* rows (or columns, since M_F is symmetric) of M_F has rank n 1.



Preliminaries 000000●	Matrix structure	Orbit restrictions	Submatrix method 00	Computational searches

Following Corollary 5 from [1], we get that function

$$F(x) = \sum_{0 \le i < j \le n-1} a_{i,j} x^{2^i + 2^j}, \ a_{i,j} \in \mathbb{F}_{2^n}$$
(2)

is APN if and only if its derivative matrix M_F is QAM.



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• Let $F(x) = \sum_{0 \le i < j \le n-1} a_{i,j} x^{2^i+2^j}$ with coefficients $a_{i,j} \in \mathbb{F}_{2^m}$ in some subfield \mathbb{F}_{2^m} of \mathbb{F}_{2^n}



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Let F(x) = ∑_{0≤i<j≤n-1} a_{i,j}x^{2ⁱ+2^j} with coefficients a_{i,j} ∈ 𝔽_{2^m} in some subfield 𝔽_{2^m} of 𝔽_{2ⁿ}
 (F(x))^{2^m} = a_i^{2^m} (xⁱ)^{2^m} = ∑_{i=0}^{2ⁿ-1} a_i (xⁱ)^{2^m} = 𝓕 (x^{2^m}),



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 (F(x))^{2^m} = a_i^{2^m} (xⁱ)^{2^m} = ∑_{i=0}^{2ⁿ-1} a_i (xⁱ)^{2^m} = F (x^{2^m}),
 (Δ_aF(x))^{2^m} = F(x+a)^{2^m} + F(x)^{2^m} + F(a)^{2^m} = Δ_{a^{2^m}}F (x^{2^m}),





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$$\begin{bmatrix} 0 & \Delta F(b_1, b_2) & \dots & \dots & \Delta F(b_1, b_n) \\ \Delta F(b_1, b_2) & 0 & \ddots & \dots & \Delta F(b_2, b_n) \\ \vdots & \ddots & \ddots & (\Delta F(b_1, b_2))^{2^m} & \vdots \\ \vdots & \ddots & (\Delta F(b_1, b_2))^{2^m} & 0 & \vdots \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ \Delta F(b_1, b_n) & \Delta F(b_2, b_n) & \dots & \dots & 0 \end{bmatrix}$$

Structure of the search

where $\Omega_1, \Omega_2, \ldots, \Omega_I \in \mathbb{F}_{2^n}$ - variables.



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Structure of the search

$$M_{F} = \begin{pmatrix} 0 & \Omega_{1} & \Omega_{2} & \dots & \dots & \dots \\ \Omega_{1} & 0 & \ddots & \ddots & \dots & \dots \\ \Omega_{2} & \dots & 0 & \Omega_{1}^{2^{m}} & \Omega_{2}^{2^{m}} & \dots \\ \vdots & \vdots & \Omega_{1}^{2^{m}} & 0 & \dots & \dots \\ \vdots & \vdots & \Omega_{2}^{2^{m}} & \dots & 0 & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \ddots \end{pmatrix},$$

where $\Omega_1, \Omega_2, \ldots, \Omega_l \in \mathbb{F}_{2^n}$ - variables. A variable Ω_i is located on the *i*-th level.

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Orbit restrictions

Theorem 3 [3]

For any linear permutation / on \mathbb{F}_{2^n} and $M \in \mathbb{F}_{2^n}^{n \times n}$ s.t. $M = M_F$ then any $M' = M_{F'}$ produced by

$$M'_{i,j} = I(M_{i,j}) \text{ for all } 1 \le i,j \le n \tag{4}$$

will be $F' = I \circ F$ linearly equivalent(also EA-equivalent) to F.



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will be $F' = I \circ F$ linearly equivalent(also EA-equivalent) to F. Let \mathcal{L} be a set of all linear (n, n)-permutations $I = \sum_{i=1}^{n} \alpha_i x^{2^{i-1}}$ on \mathbb{F}_{2^n} with subfield $\alpha_i \in \mathbb{F}_{2^m}$.



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$$Orb(a, \mathcal{L}) = \{ l(a) : l \in \mathcal{L} \}.$$
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Preliminaries 0000000	Matrix structure 00	Orbit restrictions	Submatrix method	Computational searches

Orbit Restrictions

 $\mathbb{F}_{2^n} = Orb(a_1, \mathcal{L}) \cup \cdots \cup Orb(a_k, \mathcal{L}), \text{ for some } a_i \in \mathbb{F}_{2^n}, \ 1 \leq i \leq k.$



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$$M_{F'} = \begin{pmatrix} 0 & L(\Omega_1) & L(\Omega_2) & \dots & \dots & \dots \\ L(\Omega_1) & 0 & \ddots & \ddots & \dots & \dots \\ L(\Omega_2) & \dots & 0 & L(\Omega_1^{2^m}) & L(\Omega_2^{2^m}) & \dots \\ \vdots & \vdots & L(\Omega_1^{2^m}) & 0 & \dots & \dots \\ \vdots & \vdots & L(\Omega_2^{m}) & \dots & 0 & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \ddots \end{pmatrix},$$

where
$$L(\Omega_i^{2^{m*j}}) = (L(\Omega_i))^{2^{m*j}}, j \in \{1, \ldots, n/m-1\}$$
 for any variable $\Omega_i, 1 \leq i \leq l$.



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 $\mathbb{F}_{2^n} = Orb(A, \mathcal{L}) \cup \ldots, \ A \in \mathbb{F}_{2^n}.$



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$$M_{F} = \begin{pmatrix} 0 & A & \Omega_{2} & \dots & \dots \\ A & 0 & \ddots & \ddots & \dots & \dots \\ \Omega_{2} & \dots & 0 & A^{2^{m}} & \Omega_{2}^{2^{m}} & \dots \\ \vdots & \vdots & A^{2^{m}} & 0 & \dots & \dots \\ \vdots & \vdots & \Omega_{2}^{2^{m}} & \dots & 0 & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \ddots \end{pmatrix}.$$

 $Orb_A(\Omega_2, \mathcal{L}) = \{ I(\Omega_2) : I \in \mathcal{L} \mid I(A) = A \}.$



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 $S = \{0, 0, ...\}$

$$Orb_{\mathcal{S}}(\Omega_k,\mathcal{L}) = \{ l(\Omega_k) : l \in \mathcal{L} \mid \forall X \in \mathcal{S} : l(X) = X \}.$$







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- Let $M \in \mathbb{F}_{2^n}^{n \times n}$ be a derivative matrix.
- *M* is QAM if and only if every submatrix $S \in \mathbb{F}_{2^n}^{p \times q}$, $1 \le p, q \le n$ of *M* is **proper**.



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- ► S proper if every nonzero linear combinations of the p rows has rank at least q - 1.



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After considering F' = F ∘ L, where L = a_jx^{2'}, a_j ∈ 𝔽_{2^m}, we could eliminate the number of submatrices for this test.



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Preliminaries 0000000	Matrix structure	Orbit restrictions	Submatrix method	Computational searches •000000
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Preliminaries 0000000	Matrix structure 00	Orbit restrictions	Submatrix method	Computational searches
(8.2)				

- F(x) over \mathbb{F}_{2^8} with coefficients in \mathbb{F}_{2^2} .
- 4⁸ = 65536 linear permutations with coefficients in the subfield were constructed.



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Preliminaries 0000000	Matrix structure 00	Orbit restrictions	Submatrix method	Computational searches
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(8,2)

- F(x) over \mathbb{F}_{2^8} with coefficients in \mathbb{F}_{2^2} .
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- By using these permutations, the first level of the search was partitioned into 4 orbits.



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Preliminaries 0000000	Matrix structure	Orbit restrictions	Submatrix method	Computational searches
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(8,2)

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$$1 \quad a \quad a^7 \quad a^{17}$$

Table: The first level



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Preliminaries 0000000	Matrix structure 00	Orbit restrictions	Submatrix method	Computational searches
(8.2)				

- F(x) over \mathbb{F}_{2^8} with coefficients in \mathbb{F}_{2^2} .
- 4⁸ = 65536 linear permutations with coefficients in the subfield were constructed.
- The number of variables = levels in this dimension is 8.
- By using these permutations, the first level of the search was partitioned into 4 orbit representatives.

1	а	a ⁷	a ¹⁷
$\#\{\Omega_2\}_i = 8$	$\#\{\Omega_2\}_i=30$	$\#\{\Omega_2\}_i = 22$	$\#\{\Omega_2\}_i = 14$
$Orb_1\Omega_2$	$Orb_a\Omega_2$	$Orb_{a^7}\Omega_2$	$Orb_{a^{17}}\Omega_2$

Table: The second level



Preliminaries 0000000	Matrix structure 00	Orbit restrictions	Submatrix method	Computational searches
(8.2)				
(0,2)				

1	а	a ⁷	a ¹⁷
40 hours	1 month	10 days	7 days



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Preliminaries	Matrix structure 00	Orbit restrictions	Submatrix method 00	Computational searches
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1	а	a ⁷	a ¹⁷
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 196863 quadratic APN functions were found in the search, with 27 unique ortho-derivative differential spectra.



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(8,2)

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 $\blacktriangleright \ a^{85}x^{96} + a^{85}x^{72} + a^{170}x^{24} + x^{18} + a^{85}x^{12} + a^{85}x^9 + x^6 + x^3.$



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Preliminaries 0000000	Matrix structure 00	Orbit restrictions	Submatrix method	Computational searches

(8,2)

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- 0³⁸¹⁹⁶, 2²²⁰⁰⁸, 4⁴⁶⁰⁸, 6⁴⁵⁶, 8¹² its ortho-derivative differential spectra.



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Preliminaries 0000000	Matrix structure 00	Orbit restrictions	Submatrix method 00	Computational searches
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(10.2)				

- F(x) over $\mathbb{F}_{2^{10}}$ with coefficients in \mathbb{F}_{2^2} .
- 4¹⁰ = 1048576 linear permutations with coefficients in the subfield were constructed.
- The number of variables = levels in this dimension is 9.
- By using these permutations, the first level of the search was partitioned into 3 orbit representatives.

1	а	a ⁵
$\#\{\Omega_2\}_i = 5$	$\#\{\Omega_2\}_i = 33$	$\#\{\Omega_2\}_i = 50$
$Orb_1\Omega_2$	$Orb_a\Omega_2$	$Orb_{a^5}\Omega_2$



Preliminaries	Matrix structure	Orbit restrictions	Submatrix method	Computational searches
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(10,1)				

- F(x) over $\mathbb{F}_{2^{10}}$ with coefficients in \mathbb{F}_{2^1} .
- 2¹⁰ = 1024 linear permutations with coefficients in the subfield were constructed.
- The number of variables = levels in this dimension is 5.
- By using these permutations, the first level of the search was partitioned into 8 orbit representatives.

1	а	a ⁵	a^{15}	a ³³	a ⁵⁷	a ⁹⁹	a ³⁴¹
# of orbit representatives for 2 nd level after Sub-matrix Test							
0	746	1012	753	71	112	78	8

Preliminaries	Matrix structure	Orbit restrictions	Submatrix method	Computational searches
(9,3)				

- F(x) over \mathbb{F}_{2^9} with coefficients in \mathbb{F}_{2^3} .
- 8⁹ = 134217728 linear permutations with coefficients in the subfield were constructed.
- The number of variables = levels in this dimension is 12.



Preliminaries 0000000	Matrix structure 00	Orbit restrictions	Submatrix method	Computational searches
(9.3)				

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Remark

Let $a \in \mathbb{F}_{2^9}$. We categorize *a* into the following cases:

1.
$$Cat_1 = \{a : a \in \mathbb{F}_{2^9} \mid a + a^{2^3} = 0\},\$$

2. $Cat_2 = \{a : a \in \mathbb{F}_{2^9} \mid a + a^{2^3} + a^{2^6} = 0\},\$
3. $Cat_3 = \{a : a \in \mathbb{F}_{2^9} \mid a \notin Cat_1, a \notin Cat_2\}\$



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Preliminaries 0000000	Matrix structure 00	Orbit restrictions	Submatrix method	Computational searches
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3. $Cat_3 = \{a : a \in \mathbb{F}_{2^9} \mid a \notin Cat_1, a \notin Cat_2\}\$

Theorem

Let $a, b \in Cat_3$. If there exist $l(x) = \sum_{i=0}^{8} c_i x^{2^i}$, $c_i \in \mathbb{F}_{2^3}$ s.t. l(a) = b, $l(a^{2^3}) = b^{2^3}$, $l(a^{2^6}) = b^{2^6}$. Then there exist linear permutation $L \in \mathcal{L}$ s.t. L(a) = b.



17 / 18

Preliminaries 0000000	Matrix structure 00	Orbit restrictions	Submatrix method	Computational searches

Conclusions

For F(x) over 𝔽_{2ⁿ} with coefficients in 𝔽_{2^m} we run searches (n, m) for (8, 2), (10, 2), (10, 1), (9, 3).



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Preliminaries 0000000	Matrix structure 00	Orbit restrictions	Submatrix method	Computational searches

Conclusions

- For F(x) over 𝔽_{2ⁿ} with coefficients in 𝔽_{2^m} we run searches (n, m) for (8, 2), (10, 2), (10, 1), (9, 3).
- We conclude where it is feasible to get the results and improve the computational method as possible.



Preliminaries 0000000	Matrix structure	Orbit restrictions	Submatrix method	Computational searches

Conclusions

- For F(x) over 𝔽_{2ⁿ} with coefficients in 𝔽_{2^m} we run searches (n, m) for (8, 2), (10, 2), (10, 1), (9, 3).
- We conclude where it is feasible to get the results and improve the computational method as possible.
- Computational searches are still running.

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