

On quadratic APN functions

$$F(x) + \text{Tr}(x)L(x)$$

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Abstract

We first characterize how two $(n-1, m)$ functions \mathbf{f} and \mathbf{g} can be combined into an APN (n, m) -function \mathbf{F} of the form $\mathbf{F}(\mathbf{x}) = \mathbf{f}(\mathbf{x})$ and $\mathbf{F}(\mathbf{x} + \mathbf{e}_0) = \mathbf{g}(\mathbf{x})$ for $\mathbf{x} \in \mathbb{F}_2^{n-1}$ with $\mathbf{e}_0 \in \mathbb{F}_2^n \setminus \mathbb{F}_2^{n-1}$. Next we specialize this characterization to the case when \mathbf{f} is quadratic and $\mathbf{g}(\mathbf{x}) = \mathbf{f}(\mathbf{x}) + \mathbf{L}(\mathbf{x})$ for some linearized polynomial \mathbf{L} . Lastly for a quadratic APN (n, n) -function \mathbf{F} and a linearized polynomial \mathbf{L} , we give a characterization of APN-ness for (n, n) -function $\mathbf{F}(\mathbf{x}) + \text{Tr}(\mathbf{x})\mathbf{L}(\mathbf{x})$. With some computational experiments, we see that CCZ-inequivalent APN functions $\mathbf{F}(\mathbf{x}) + \text{Tr}(\mathbf{x})\mathbf{L}(\mathbf{x})$ can be obtained from \mathbf{F} using this construction.

1 Preliminaries

Let \mathbb{F}_2 be the binary field, and n, m positive integers. A function $F : \mathbb{F}_2^n \rightarrow \mathbb{F}_2^m$ is called an almost perfect nonlinear (APN) function if the cardinality $|\{x \mid F(x+a) + F(x) = b\}|$ is less than or equal to 2 for any nonzero $a \in \mathbb{F}_2^n$ and for any $b \in \mathbb{F}_2^m$. APN functions have been studied for many years because of their applications in cryptography. See [1], [2] or [5] for known APN functions. We call a function F *quadratic* if $F(x+y) + F(x) + F(y) + F(0)$ is \mathbb{F}_2 -bilinear. Two functions F_1 and F_2 from \mathbb{F}_2^n to \mathbb{F}_2^m are called *CCZ-equivalent* if the graphs $G_{F_1} := \{(x, F_1(x)) \mid x \in \mathbb{F}_2^n\}$ and $G_{F_2} := \{(x, F_2(x)) \mid x \in \mathbb{F}_2^n\}$ in $\mathbb{F}_2^n \oplus \mathbb{F}_2^m$ are affine equivalent, that is, if there exists an \mathbb{F}_2 -linear isomorphism $l \in GL_2(\mathbb{F}_2^n \oplus \mathbb{F}_2^m)$ and an element $v \in \mathbb{F}_2^n \oplus \mathbb{F}_2^m$ such that $l(G_{F_1}) + v = G_{F_2}$. The Γ -rank of a function $F : \mathbb{F}_2^n \rightarrow \mathbb{F}_2^m$ is the rank of the incidence matrix over

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\mathbb{F}_2 of the incidence structure $\{\mathcal{P}, \mathcal{B}, I\}$, where $\mathcal{P} = \mathbb{F}_2^n \oplus \mathbb{F}_2^m$, $\mathcal{B} = \mathbb{F}_2^n \oplus \mathbb{F}_2^m$ and $(a, b)I(u, v)$ for $(a, b) \in \mathcal{P}$ and $(u, v) \in \mathcal{B}$ if and only if $F(a + u) = b + v$. We know that if two functions F_1 and F_2 from \mathbb{F}_2^n to \mathbb{F}_2^m are CCZ-equivalent, then they have the same Γ -rank (see [3]). Let \mathbb{F}_{2^n} be the finite field of 2^n elements. We sometimes identify \mathbb{F}_{2^n} with \mathbb{F}_2^n as an \mathbb{F}_2 -vector space. We denote the set $\mathbb{F}_{2^n} \setminus \{0\}$ by $\mathbb{F}_{2^n}^\times$ and $\mathbb{F}_2^n \setminus \{0\}$ by $(\mathbb{F}_2^n)^\times$. For finite fields $K \supset F$ of characteristic 2, we denote the trace function from K to F by Tr_F^K . We denote $\text{Tr}_{\mathbb{F}_2^K}$ by Tr and call it the absolute trace of K .

For a function F on \mathbb{F}_{2^n} , the value at $a \in \mathbb{F}_{2^n}$ of the Walsh transformation of the Boolean function $\mathbb{F}_{2^n} \ni x \mapsto \text{Tr}(bF(x)) \in \mathbb{F}_2$ for $b \in \mathbb{F}_{2^n}^\times$ is defined by

$$W_F(a, b) = \sum_{x \in \mathbb{F}_{2^n}} (-1)^{\text{Tr}(bF(x) + ax)}.$$

The Walsh spectrum of F is defined by $\mathcal{W}_F = \{W_F(a, b) \mid a \in \mathbb{F}_{2^n}, b \in \mathbb{F}_{2^n}^\times\}$. For a quadratic APN function F on \mathbb{F}_{2^n} , it is known that $W_F \in \{0, \pm 2^{(n+1)/2}\}$ if n is odd. For the case n is even, it is said that a quadratic APN function F has the classical Walsh spectrum if $\mathcal{W}_F = \{0, \pm 2^{n/2}, \pm 2^{(n+2)/2}\}$, and F has the non-classical Walsh spectrum if otherwise (see [4]).

2 A condition to have an APN function F from \mathbb{F}_2^n to \mathbb{F}_2^m using APN functions f, g from \mathbb{F}_2^{n-1} to \mathbb{F}_2^m

Let f, g be functions from \mathbb{F}_2^{n-1} to \mathbb{F}_2^m . We regard $\mathbb{F}_2^{n-1} \subset \mathbb{F}_2^n$ as an F_2 -linear subspace. Let $e_0 \in \mathbb{F}_2^n$ with $e_0 \notin \mathbb{F}_2^{n-1}$ and $\mathbb{F}_2^{n-1} + e_0 := \{x + e_0 \mid x \in \mathbb{F}_2^{n-1}\}$. Then $\mathbb{F}_2^n = \mathbb{F}_2^{n-1} \cup (\mathbb{F}_2^{n-1} + e_0)$. We want to have an APN function F from $\mathbb{F}_2^n = \mathbb{F}_2^{n-1} \cup (\mathbb{F}_2^{n-1} + e_0)$ to \mathbb{F}_2^m defined by $F(x) = f(x)$ and $F(x + e_0) = g(x)$ for $x \in \mathbb{F}_2^{n-1}$.

Proposition 1 *F defined above is an APN function if and only if*

- (1) f and g are APN functions from \mathbb{F}_2^{n-1} to \mathbb{F}_2^m ,
- (2) $f(x + a) + f(x) \neq g(y + a) + g(y)$ for any $x, y \in \mathbb{F}_2^{n-1}$ and for any nonzero $a \in \mathbb{F}_2^{n-1}$, and
- (3) $G_a : \mathbb{F}_2^{n-1} \ni x \mapsto f(x + a) + g(x) \in \mathbb{F}_2^m$ are one-to-one mappings for any $a \in \mathbb{F}_2^{n-1}$.

Proof Recall that F is an APN function if and only if, for any nonzero $A \in \mathbb{F}_2^n$ and for $X, Y \in \mathbb{F}_2^n$, $F(X + A) + F(X) = F(Y + A) + F(Y)$ implies $X = Y$ or $X = Y + A$.

Firstly assume that F is an APN function, and we will see that f and g must satisfy the conditions (1), (2) and (3).

Let $A = a \in (\mathbb{F}_2^{n-1})^\times$. For any $Y = y \in \mathbb{F}_2^{n-1}$, we must have $X = y \in \mathbb{F}_2^{n-1}$ or $X = y + a \in \mathbb{F}_2^{n-1}$ from $F(X + a) + F(X) = F(y + a) + F(y)$. Since $X \in \mathbb{F}_2^{n-1}$, we

have $f(X+a) + f(X) = f(y+a) + f(y)$ from $F(X+a) + F(X) = F(y+a) + F(y)$. Thus f must be an APN function. Next, for any $Y = y + e_0$ with $y \in \mathbb{F}_2^{n-1}$ we must have $X = y + e_0$ or $X = y + a + e_0$ from $F(X+a) + F(X) = F(y+e_0+a) + F(y+e_0)$. Since $X = x + e_0$ for some $x \in \mathbb{F}_2^{n-1}$, we have $g(x+a) + g(x) = g(y+a) + g(y)$ from $F(X+a) + F(X) = F(y+e_0+a) + F(y+e_0)$. Hence g must be an APN function. Thus the condition (1) must be satisfied.

Let $A = a \in (\mathbb{F}_2^{n-1})^\times$. For any $Y = y \in \mathbb{F}_2^{n-1}$, since $X = y$ or $X = y + a$, $F(X+a) + F(X) = F(y+a) + F(y)$ does not have a solution $X = x + e_0$ for $x \in \mathbb{F}_2^{n-1}$. Thus $F(x+e_0+a) + F(x+e_0) \neq F(y+a) + F(y)$ for any $x, y \in \mathbb{F}_2^{n-1}$, therefore we must have $g(x+a) + g(x) \neq f(y+a) + f(y)$ for any $x, y \in \mathbb{F}_2^{n-1}$. Thus the condition (2) must be satisfied.

Let $A = a + e_0$ with $a \in \mathbb{F}_2^{n-1}$ and $Y = y \in \mathbb{F}_2^{n-1}$. We have $X = y \in \mathbb{F}_2^{n-1}$ or $X = y + a + e_0$ with $y + a \in \mathbb{F}_2^{n-1}$. For $X \in \mathbb{F}_2^{n-1}$, we have $g(X+a) + f(X) = g(y+a) + f(y)$ from $F(X+a+e_0) + F(X) = F(y+a+e_0) + F(y)$, hence $g(X+a) + f(X) = g(y+a) + f(y)$ must have only one solution $X = y$ for any $y, a \in \mathbb{F}_2^{n-1}$. For $X \notin \mathbb{F}_2^{n-1}$, we have $f(X+a) + g(X) = g(y+a) + f(y)$ from $F(X+a) + F(X+e_0) = F(y+a+e_0) + F(y)$, hence $f(X+a) + g(X) = g(y+a) + f(y)$ must have only one solution $X = y + a$. Thus we see that the condition (3) must be satisfied.

Conversely, let us assume the conditions (1), (2) and (3). Assume $F(X+A) + F(X) = F(Y+A) + F(Y)$ with $A \neq 0$. We will prove that $X = Y$ or $X = Y + A$. We divide the case into the following four cases (i) $A = a \in (\mathbb{F}_2^{n-1})^\times$ and $Y = y \in \mathbb{F}_2^{n-1}$, (ii) $A = a \in (\mathbb{F}_2^{n-1})^\times$ and $Y = y + e_0$ with $y \in \mathbb{F}_2^{n-1}$, (iii) $A = a + e_0$ with $a \in \mathbb{F}_2^{n-1}$ and $Y = y$ with $y \in \mathbb{F}_2^{n-1}$, and (iv) $A = a + e_0$ with $a \in \mathbb{F}_2^{n-1}$ and $Y = y + e_0$ with $y \in \mathbb{F}_2^{n-1}$.

Firstly let us consider the case (i). If $X = x \in \mathbb{F}_2^{n-1}$, then we have $f(x+a) + f(x) = f(y+a) + f(y)$ hence $x = y$ or $x = y + a$ by (1). Let $X = x + e_0$ with $x \in \mathbb{F}_2^{n-1}$, then we have $g(x+a) + g(x) = f(y+a) + f(y)$ which has no solution by (2). Therefore, $X = Y$ or $X = Y + A$ in case (i).

Next, we consider the case (ii). Assume $X = x \in \mathbb{F}_2^{n-1}$, then we have $f(x+a) + f(x) = g(y+a) + g(y)$ which has no solution by (2). If $X = x + e_0$ with $x \in \mathbb{F}_2^{n-1}$, then we have $g(x+a) + g(x) = g(y+a) + g(y)$, hence $x + e_0 = y + e_0$ or $x + e_0 = y + e_0 + a$ by (1). Thus we have $X = Y$ or $X = Y + A$ in case (ii).

Let us consider the case (iii). If $X = x \in \mathbb{F}_2^{n-1}$, then we have $g(x+a) + f(x) = g(y+a) + f(y)$. Since $G_a : x + a \mapsto f(x) + g(x+a)$ is a one-to-one mapping by (3), we have $x = y$. If $X = x + e_0$ with $x \in \mathbb{F}_2^{n-1}$, then we have $f(x+a) + g(x) = g(y+a) + f(y)$. By the same reason as above, we have $x + e_0 = y + (a + e_0)$. Thus we have $X = Y$ or $X = Y + A$ in case (iii).

Lastly we consider the case (iv). If $X = x \in \mathbb{F}_2^{n-1}$, then we have $g(x+a) + f(x) = f(y+a) + g(y)$. Since $G_a : x \mapsto f(x+a) + g(x)$ is a one-to-one mapping by (3), we have $x = (y + e_0) + (a + e_0)$. If $X = x + e_0$ with $x \in \mathbb{F}_2^{n-1}$, then we have $f(x+a) + g(x) = f(y+a) + g(y)$. By the same reason as above, we have $x + e_0 = y + e_0$. Thus we also have $X = Y$ or $X = Y + A$ in case (iv).

Hence F must be an APN function under the conditions (1), (2) and (3). \square

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3 The case f is a quadratic APN function and $g(x) = f(x) + L'(x)$ with L' a linear mapping

Let f be a function from \mathbb{F}_2^{n-1} to \mathbb{F}_2^m and $B_f(x, a) := f(x+a) + f(x) + f(a) + f(0)$. Recall that f is quadratic if $B_f(x, a)$ is an \mathbb{F}_2 -bilinear mapping. In this section, we consider the case that f is a quadratic APN function from \mathbb{F}_2^{n-1} to \mathbb{F}_2^m , and $g(x) = f(x) + L'(x)$ for $x \in \mathbb{F}_2^{n-1}$ with L' an \mathbb{F}_2 -linear mappings from \mathbb{F}_2^{n-1} to \mathbb{F}_2^m . We note that, if f is quadratic, $\mathbb{F}_2^{n-1} \ni x \mapsto L'(x) + B_f(x, a) \in \mathbb{F}_2^m$ are linear mappings for any $a \in \mathbb{F}_2^{n-1}$. We check the conditions (1), (2) and (3) in Proposition 1. We regard \mathbb{F}_2^{n-1} as an $(n-1)$ -dimensional subspace of \mathbb{F}_2^n .

Proposition 2 *Let f be a quadratic APN function from \mathbb{F}_2^{n-1} to \mathbb{F}_2^m , and $g(x) = f(x) + L'(x)$ with L' an \mathbb{F}_2 -linear mapping from \mathbb{F}_2^{n-1} to \mathbb{F}_2^m . Let F be a function from \mathbb{F}_2^n to \mathbb{F}_2^m defined in Section 2, that is, $F(x) := f(x)$ and $F(x + e_0) := f(x) + L'(x)$ for some fixed $e_0 \in \mathbb{F}_2^n \setminus \mathbb{F}_2^{n-1}$ for $x \in \mathbb{F}_2^{n-1}$. Then F is an APN function if and only if $\mathbb{F}_2^{n-1} \ni x \mapsto L'(x) + B_f(x, a) \in \mathbb{F}_2^m$ are one-to-one mappings for any $a \in \mathbb{F}_2^{n-1}$.*

Proof Since f and $g = f + L'$ are APN functions, the condition (1) is satisfied. The condition (2) implies $f(x+a) + f(x) \neq f(y+a) + f(y) + L'(a)$ for any $x, y \in \mathbb{F}_2^{n-1}$ if $a \neq 0$, that is, $L'(a) + (f(x+a) + f(x)) + (f(y+a) + f(y)) \neq 0$ for any $x, y \in \mathbb{F}_2^{n-1}$ if $a \neq 0$, which means $L'(a) + B_f(a, x+y) \neq 0$ if $a \neq 0$, $a \in \mathbb{F}_2^{n-1}$. The condition (3) implies $G_a : \mathbb{F}_2^{n-1} \ni x \mapsto f(x+a) + g(x) = L'(x) + (f(x+a) + f(x)) \in \mathbb{F}_2^m$ are one-to-one mappings for any $a \in \mathbb{F}_2^{n-1}$, that is, $\mathbb{F}_2^{n-1} \ni x \mapsto L'(x) + B_f(x, a) + f(a) \in \mathbb{F}_2^m$ are one-to-one mappings for any $a \in \mathbb{F}_2^{n-1}$. Thus we see that the conditions (1), (2) and (3) in Proposition 1 are satisfied if and only if $\mathbb{F}_2^{n-1} \ni x \mapsto L'(x) + B_f(x, a) \in \mathbb{F}_2^m$ are one-to-one mappings for any $a \in \mathbb{F}_2^{n-1}$. \square

4 $F(x) + \text{Tr}(x)L(x)$ for a quadratic APN function F on \mathbb{F}_{2^n}

Let $T_0 := \{x \in \mathbb{F}_{2^n} \mid \text{Tr}(x) = 0\}$ and $e_0 \in \mathbb{F}_{2^n}$ with $\text{Tr}(e_0) = 1$. Let F be a quadratic APN function on \mathbb{F}_{2^n} and $B_F(x, a) := F(x+a) + F(x) + F(a) + F(0)$ for $x, a \in \mathbb{F}_{2^n}$. Let L be an \mathbb{F}_2 -linear mapping on \mathbb{F}_{2^n} .

Theorem 3 *Let F be a quadratic APN function on \mathbb{F}_{2^n} and L an \mathbb{F}_2 -linear mapping on \mathbb{F}_{2^n} . Let $e_0 \in \mathbb{F}_{2^n}$ with $\text{Tr}(e_0) = 1$. Then, $F(x) + \text{Tr}(x)L(x)$ is a quadratic APN function on \mathbb{F}_{2^n} if and only if $L_a : T_0 \ni x \mapsto L(x) + B_F(x, a + e_0) \in \mathbb{F}_{2^n}$ are one-to-one mappings from T_0 to \mathbb{F}_{2^n} for any $a \in T_0$. (Hence, $F(x) + \text{Tr}(x)L(x)$ is a quadratic APN function on \mathbb{F}_{2^n} if, and only if, $L_a(x) = 0$ implies $x = 0$ for any $a \in T_0$).*

Proof Let $f := F|_{T_0}$ be the restriction of F to T_0 ; f is a quadratic APN function from T_0 to \mathbb{F}_{2^n} . For $x \in T_0$, we have $F(x) + \text{Tr}(x)L(x) = f(x)$ and $F(x + e_0) +$

$\text{Tr}(x + e_0)L(x + e_0) = f(x) + L(x) + B_F(x, e_0) + L(e_0) + F(e_0)$. Let G be a function on \mathbb{F}_{2^n} defined by $G(x) := f(x)$ and $G(x + e_0) := f(x) + L(x) + B_F(x, e_0)$ for $x \in T_0$, then $G(x) = F(x) + \text{Tr}(x)(L(x) + L(e_0) + F(e_0))$ for $x \in \mathbb{F}_{2^n}$, which is CCZ equivalent to $F(x) + \text{Tr}(x)L(x)$. By Proposition 2, G is an APN function if and only if $T_0 \ni x \mapsto L(x) + B_F(x, e_0) + B_F(x, a) \in \mathbb{F}_{2^n}$ are one-to-one mappings for any $a \in T_0$. Thus $F(x) + \text{Tr}(x)L(x)$ is a quadratic APN function on \mathbb{F}_{2^n} if and only if $L_a : T_0 \ni x \mapsto L(x) + B_F(x, a + e_0) \in \mathbb{F}_{2^n}$ are one-to-one mappings from T_0 to \mathbb{F}_{2^n} for any $a \in T_0$. \square

Let e_0 be some fixed element of \mathbb{F}_{2^n} with $\text{Tr}(e_0) = 1$. Using a computer, for linear mappings L on \mathbb{F}_{2^n} such that $L_a : T_0 \ni x \mapsto L(x) + B(x, a + e_0) \in \mathbb{F}_{2^n}$ are one-to-one mappings from T_0 to \mathbb{F}_{2^n} for any $a \in T_0$, we have 448 L 's with $L(e_0) = 0$ for $F(x) = x^3$ on \mathbb{F}_{2^4} , 4608 L 's with $L(e_0) = 0$ for $F(x) = x^3$ on \mathbb{F}_{2^5} , and many (about 40,000) L 's with $L(e_0) = 0$ for $F(x) = x^3$ on \mathbb{F}_{2^6} .

Example 1 Let $F(x) = x^3$ on \mathbb{F}_{2^6} . The Γ -rank of F is 1102. Using a computer, we see that there are linear mappings L satisfying the conditions in Theorem 3 such that the Γ -ranks of $F(x) + \text{Tr}(x)L(x)$ are 1144, 1146, 1158, 1166, 1168, 1170, 1172 and 1174. We also see that $F(x) + \text{Tr}(x)L(x)$ with $L(x) = \alpha^{42}x + \alpha^{19}x^2 + \alpha^{51}x^2^2 + \alpha^{59}x^2^3 + \alpha^{26}x^2^4 + \alpha^{38}x^2^5$, where α is a primitive element of \mathbb{F}_{2^6} , has non-classical Walsh spectrum $\mathcal{W}_F = \{0, \pm 8, \pm 16, \pm 32\}$ with the Γ -rank 1170. Since $F(x) + \text{Tr}(x)L(x)$ with $L(x) = \alpha^{42}x + \alpha^{47}x^2 + \alpha^{35}x^2^2 + \alpha^{54}x^2^3 + \alpha^{23}x^2^4 + \alpha^{27}x^2^5$ has classical Walsh spectrum $\mathcal{W}_F = \{0, \pm 8, \pm 16\}$ with the Γ -rank 1170, we see that there are inequivalent APN functions $F(x) + \text{Tr}(x)L(x)$ with the same Γ -rank.

Let $F(x) = x^3$ on \mathbb{F}_{2^7} . The Γ -rank of F is 3610. Using a computer, we find that the linear mapping $L(x) := x + x^{2^3} + x^{2^5} + x^{2^6}$ satisfies the conditions in Theorem 3 and the Γ -rank of $F(x) + \text{Tr}(x)L(x)$ is 4048.

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