# **Cryptanalysis – a Never-Ending Story**

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Ínría

## **Big data = big security breaches**



https://www.databreachtoday.com/marriott-mega-breach-victim-count-drops-to-383-million-a-11916

## Sensitive data

# Healthcare under Attack: What Happens to Stolen Medical Records?

30 czerwca 2016



According to **reports** from news portal Deep Dot Web, 689,621 patient records are being sold by a hacker operating in TheRealDeal, a deep web marketplace **known** for peddling stolen data, codes and **zero-day** software exploits. The hacker told the news site that he used an exploit in how the organizations utilize remote desktop protocol (RDP), adding that it is a specific security flaw with precise conditions needed for it to be triggered.

The hacker, who goes by the handle "thedarkoverlord," is offering a purportedly one-off copy of the stolen data. The data has been broken into databases, with prices ranging

🖬 f Powiązane arty Unsecure Pagers in Vancouver Expose S Patient Data: What Means for Enterprise Recognizing Enterpr Mission-Critical Ass License Plates, Phote Passwords and More Two Separate Bread Ransomware Hits Co > Offices, Knocks The Channel Offline Android Malware Ca SimBad Adware and Operation Sheep Re Installed 250 Million Najnowsze art

#### https://www.trendmicro.com/vinfo/pl/security/news/cyber-attacks/

#### ${\tt health care-under-attack-stolen-medical-records}$

## **Attacks**

## In most attacks, cryptography is bypassed.

"I am not aware of any major world-class security system employing cryptography in which the hackers penetrated the system by actually going through the cryptanalysis." [Adi Shamir 2002]

## **Attacks**



https:

//afroginthefjord.com/2014/07/12/why-the-french-are-so-arrogant-and-why-norwegians-arent/

However...

- biases in RC4 [AlFardam et al. 13]
- Logjam [Adrian et al. 15]: weak Diffie-Helman
- Sloth [Bhargavan, Leurent 16]: collisions in MD5

## FLAME spy malware



The malware appears as a Windows Update security patch, with a valid certificate.

The fraudulent signature has been forged with a chosen-prefix collision attack against MD5 [Stevens 07]

## **Attacks against MIFARE**

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Home > News > IT Vendors > Questions raised about Oyster card security

# Questions raised about Oyster card security

Its RFID chip is cracked by researchers

Network World and Computerworld UK staff March 7, 2008

re Smartcards with encrypted RFID chips, including London's Oyster fare card, might not be as secure as previously thought.

Can we trust cryptographers?

## Kerckhoffs' principles (1883)

The system must not require secrecy and can be stolen by the enemy without causing trouble.

The Administration must absolutely renounce secret methods, and must establish in principle that it will only accept a process that can be taught publicly in our military schools, that our students will be free to to communicate to whomever they want.

A cipher

K is a secret key.



## Perfect secrecy [Shannon 49]

"For all ciphertexts, the *a posteriori* probabilities for the various messages are equal to the *a priori* probabilities independently of the values of these."

Intercepting the message has given the cryptanalyst no information.

For any pair (m,c),

$$\Pr[M = m | \mathbf{C} = \mathbf{c}] = \Pr[M = m].$$

Equivalently,

$$\Pr[C=c|M=m]=\Pr[C=c].$$

There is at least one key transforming any plaintext m into any of the ciphertexts c.

 $\longrightarrow$  nb of keys  $\ge$  nb of possible plaintexts.

## Vernam cipher [Vernam 1926]

plaintext	u	n	b	r	е	а	k	а	b	I	е
	20	13	1	17	4	0	10	0	1	11	4
secret key	n	W	Ι	r	b	b	m	q	b	h	С
	13	22	11	17	1	1	12	16	1	7	2
ciphertext	7	9	12	8	5	1	22	16	2	18	6
	h	j	m	i	f	b	W	q	С	S	g

The Signal Corps tested the secrecy of communications handled by the system and tried it out between New-York and Washington. This trial proved that the system could be successfully used to send messages secretly.

## Vernam cipher [Vernam 1926]

```
plaintextu n b r e a k a b l esecret keyn w l r b b m q b h cciphertexth j m i f b w q c s g
```

```
plaintext in formation
secret key zwhuopwxuet
ciphertext hjmifbwqcsg
```

The plaintext and the ciphertext are statistically independent.

$$\Pr[C=c|M=m]=\Pr[C=c].$$

## **Practical secrecy**

measure of the amount of work required to break the system

## **Breaking:**

- recover the plaintext from the ciphertext;
- recover the key from the knowledge of some plaintext-ciphertext pairs.

## Amount of work:

- time;
- memory;
- data.

## Paranoia?

## **Cryptanalysis of the full Spritz** [Banik, Isobe 16]:

"We need approximatively  $2^{1247}$  assignments to recover the internal state."

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## **Cryptanalysis of the full Spritz** [Banik, Isobe 16]:

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$$2^{1247}\simeq 10^{375}\gg (\# ext{ atoms in the universe })^4$$

What is cryptanalysis about?

A good primitive must behave as a function chosen at random from the set of all functions with the same characteristics.

## **Symmetric Encryption Schemes**

## For encrypting messages of an arbitrary length:

- use a transformation operating on *n*-bit blocks (block cipher)
- chain the blocks with a mode of operation (CBC, CTR...)



**Typical block size:** 

$$n\in\{128,64\}$$

## Ideal block cipher

A block cipher operating on n-bit messages with a k-bit key

### $\iff$

 $2^k$  permutations of the set of n-bit messages, randomly selected among the  $2^n!$  possible ones.

## Ideal block cipher

A block cipher operating on n-bit messages with a k-bit key

 $2^k$  permutations of the set of n-bit messages, randomly selected among the  $2^n!$  possible ones.

**Requirement.** These permutations P cannot be distinguished from a randomly selected permutation from the knowledge of some pairs (x, P(x)).

## **Ideal behavior**

A block cipher operating on n-bit messages with a k-bit key

 $\iff$ 

 $2^k$  permutations of the set of *n*-bit messages, randomly selected among the  $2^n!$  possible ones.

No attack better than brute-force: There is no attack having a complexity significantly less than the cost of  $2^k$  evaluations of the cipher.

## **Differential cryptanalysis** [Biham-Shamir 90]



For a random permutation  $\pi$  of n-bit messages, for any nonzero a and b,

$$\Pr_X[\pi(X+\boldsymbol{a})-\pi(X)=\boldsymbol{b}] = \frac{1}{2^n-1}$$

## **Broken?**

A good primitive must behave as a function chosen at random from the set of all functions with the same characteristics.

## Cryptanalysis of the full Spritz [Banik, Isobe 16]:

"We need approximatively  $2^{1247}$  assignments to recover the internal state."

## **Specifications of Spritz** [Rivest, Schulz 15]:

Spritz generates a pseudo-random sequence from a secret state, chosen out of  $2^{1730}$  possibilities.

The internal state can be recovered with  $2^{1247}$  trials ightarrow much better than brute-force.

**Practical relevance?** 

But this is not relevant in our applications...

Finding collisions is not an issue in key-exchange protocols.

Sloth attack against TLS [Bhargavan, Leurent 16]: exploits collisions in MD5!

## But these attacks are not practical...

Attacks reveal unexpected weaknesses.

## But these attacks are not practical...

Attacks reveal unexpected weaknesses.

Attacks always get better; they never get worse.

If cryptographers say that an algorithm is broken, don't use it!

Is there any difference between

- **AES** (NIST FIPS 197)
- **Crypto-1** (MIFARE Classic encryption)
- **Dual-EC-DRBG** (NIST SP 800-90A)

AES has been standardized after an open competition (1997-2001)

The foundation of trust

## Hash function competition (SHA-3)

## Oct 2008 submission deadline

 $\longrightarrow$  64 candidates received by the NIST

- Dec 2008 51 candidates in the 1st round
- Feb 2009 1st SHA-3 conference

## Let's start the struggle!

Abacus	Neil Sholer	in round 1	2nd-preimage	
ARIRANG	Jongin Lim	in round 1		
AURORA Masahiro Fuji		in round 1	2nd preimage	
Blender	Colin Bradbury	in round 1 collision, preimage		near-collision
Boole	Greg Rose	in round 1	collision	
Cheetah	Dmitry Khovratovich	in round 1		length- extension
CHI	Phillip Hawkes	in round 1		
CRUNCH	Jacques Patarin	in round 1		length- extension
DCH	David A. Wilson	in round 1	collision	
Dynamic SHA	Xu Zijie	in round 1	collision	length- extension

http://ehash.iaik.tugraz.at/wiki/The\_SHA-3\_Zoo

## Hash function competition (SHA-3)

## Oct 2008 submission deadline

- $\longrightarrow$  64 candidates received by the NIST
- Dec 2008 51 candidates in the 1st round
- Feb 2009 1st SHA-3 conference
- July 2009 14 candidates in the 2nd round
- Aug 2010 2nd SHA-3 conference
- Dec 2010 5 finalists
- Mar 2012 3rd SHA-3 conference
- Oct 2012 winner announced (Keccak)

## **Prize for the best cryptanalysis**

### Third cryptanalysis prize

30 September 2009

We announce the third prize for the most interesting cryptanalysis of KECCAK. The results must be publicly available on an URL that is sent to keccak *-at-* noekeon *-dot-* org **before December 5, 2009** at 23:59 GMT+1 (i.e., before Sinterklaas or Saint Nicolas).

The third prize consists of beer, like the first one. This time we offer Lambic beers that according to myth can only be brewed in the surroundings of Brussels thanks to wild yeast and mysterious bacteria that would not occur anywhere else. Anyway, the prize is a case with 24 (the new number of rounds in KECCAK-*f*) bottles of Lambic-based beers from breweries such as Cantillon, Girardin, and 3 Fonteinen.

http://keccak.noekeon.org/third\_party.html

## **Prize for the best cryptanalysis**

[Boura, Canteaut 2011]: distinguisher on the inner permutation of Keccak with complexity  $2^{1575}$  (instead of  $2^{1600}$ ).



## **Iterated block ciphers**



where F is a keyed permutation of  $\{0,1\}^n$ .

## AES-128 (10 rounds) [Daemen, Rijmen 98]:

5 rounds 2 <sup>46</sup>		Daemen, Rijmen 1998
5 rounds	2 <sup>24</sup>	Bar-On et al. 2018
6 rounds	2 <sup>71</sup>	Daemen, Rijmen 1998
6 rounds	2 <sup>48</sup>	Ferguson et al. 2000
7 rounds	$\simeq 2^{128}$	Gilbert, Minier 2000
7 rounds	$2^{117}$	Lu, Dunkelman, Keller, Kim 2008
7 rounds	$2^{110}$	Mala et al. 2010
7 rounds	2 <sup>99</sup>	Derbez, Fouque, Jean 2013

 $\rightarrow$  a never-ending evaluation of the security margin is needed "Stay critical!" [Daemen 11]

## **Cryptanalysis: foundation of trust**

No public analysis, no trust

## **Examples:**

- Crypto-1 (Mifare): proprietary design
- Dual-EC-DRBG: backdoor
- Simon, Speck [NSA 2015]: no design rationale
- Streebog, Kuznyechik [FSB 2015]: structure that cannot possibly be the outcome of a random generation process, contrary to the claims of the designers [Perrin 19]

New targets

## **New functionnalities**



## **New threats**

# **Forbes**



## SEP 4, 2015 | FORBES Forbes: Quantum Computing: From Theory To Reality

#### Author: Jason Bloomberg

"The word quantum often portends New Age mumbo-jumbo, in spite of the fact that quantum mechanics underlies many of today's most important technologies, including lasers and the semiconductors found in every computer chip.

## New implementation constraints



## **NIST Lightweight Competition**

https://csrc.nist.gov/Projects/lightweight-cryptography

Feb 2019 submission deadline

 → 57 candidates received by the NIST
 Aug 2019 32 candidates in the 2nd round
 Nov 2019 NIST lightweight workshop

 Sept 2020? finalists

## **2nd-round lightweight candidates**

ACE DryGASCON ForkAE Grain-128AEAD KNOT ORANGE Pyjamask **S**aturnin SPIX Subterranean 2.0 WAGE

ASCON Elephant **GIFT-COFB HyENA** LOTUS/LOCUS-AEAD Oribatida Romulus SKINNY-AEAD SpoC SUNDAE-GIFT Xoodyak

COMET ESTATE Gimli ISAP mixFeed PHOTON-Beetle SAEAES SPARKLE Spook TinyJAMBU Conclusion

# **Public analysis is the only reliable security argument**